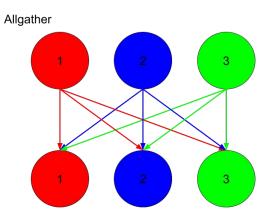
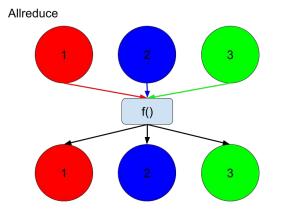
Device Memory TCP Transferring data from/to device memory efficiently

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Problem space: Large ML jobs

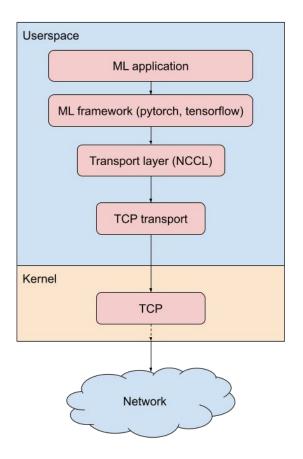
- Machine learning jobs that span many nodes.
- Data held on each nodes in **GPU** memory.
- Jobs requires efficient data transfer between GPUs on different nodes.





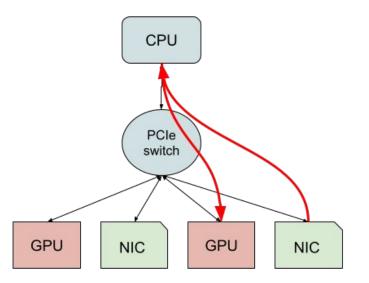
Problem space: Node zoom-in

- TCP requires a host memory bounce buffer.
 - Consumes memory bandwidth.
 - Consumes PCIe bandwidth.



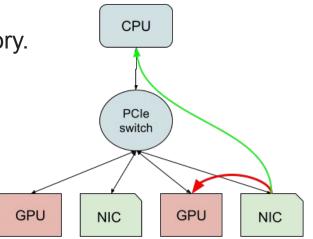
Problem space: PCle bandwidth utilization

- Data goes from NIC -> root complex (host bounce buffers) -> GPU and vice versa.
- Stresses shared PCIe links.
- Google Cloud <u>A3 VMs</u> (8 H100 GPUs)
- Nvidia DGX H100 Systems (8 H100 GPUs)



Proposed solution: Device memory TCP

- Eliminate host memory bounce buffer.
- Transfer data directly to/from device memory.
- Packet payload lands directly into device memory.
- Packet header lands into host memory.
- <u>**RFC</u>** on the mailing list.</u>



Journey of an RX packet: Device memory setup.

- Device memory abstraction of choice: dma-buf.
 - Standard in-kernel abstraction for device memory.
 - Device memory owner is an 'exporter'.
 - Device memory user is an 'importer'.
 - dma-buf APIs handle mapping/unmapping.
 - Not struct paged...
- User allocates device memory, obtains a dma-buf handle to the device memory.

Journey of an RX packet: NIC setup

- User 'binds' RX-queue to dma-buf.
 - Netdev netlink APIs
 - page_pool handles memory allocation from the dma-buf.
- Configures RSS to steer all other traffic to other queues
 - ethtool -X <if> equal 15
- Configures flow steering to steer their traffic to that queue:
 - ethtool -N <if> src-ip <ip> dst-ip <ip>... queue 15

Journey of an RX packet: page_pool

Idea based on Jakub's memory-provider <u>RFC</u>.



 Enables plugging in 'memory-providers' to the page_pool, supporting different memory types.

Journey of an RX packet: page_pool

- Dma-buf memory provider takes care of allocating PAGE_SIZE slices from the dma-buf and feeding them to the page_pool.
- But, dma-buf has no pages...
 page_pool_iovs!

<pre>struct page_pool_iov { struct dmabuf_genpool_chunk_owner *owner;</pre>
refcount_t refcount;
<pre>struct page_pool *pp; };</pre>
dma_addr_t page_pool_iov_dma_addr(<mark>const</mark> struct page_pool_iov *ppiov);
unsigned long page_pool_iov_virtual_addr(<mark>const</mark> struct page_pool_iov *ppiov);
<pre>int page_pool_iov_refcount(const struct page_pool_iov *ppiov);</pre>
<pre>void page_pool_iov_get_many(struct page_pool_iov *ppiov,</pre>
<pre>void page_pool_iov_put_many(struct page_pool_iov *ppiov,</pre>

Journey of an RX packet: nonpaged memory support

- page_pool, drivers, and skb_frag_t all use page* today.
- So much code churn... LSB pointer trick to reduce code churn.
- The LSB on page_pool_iov* is set and it's cast to page*.

#define PP DEVMEM 0x01UL static struct page *mp_dmabuf_devmem_alloc_pages(struct page_pool *pool, gfp_t gfp) ppiov = netdev_alloc_devmem(binding); ppiov = (struct page_pool_iov *)((unsigned long)ppiov | PP_DEVMEM); return (struct page *)ppiov; static inline bool page_is_page_pool_iov(const struct page *page) return (unsigned long)page & PP_DEVMEM; static inline struct page pool iov *page to page pool iov(struct page *page) if (page_is_page_pool_iov(page)) return (struct page_pool_iov *)((unsigned long)page & ~PP_DEVMEM); DEBUG_NET_WARN_ON_ONCE(true); return NULL;

Journey of an RX packet: nonpaged memory support

- page_pool does most of the heavy lifting.
- Handles any special casing for page_pool_iov
- Refcounting, dma_addr handling, pp info.
- Page-recycling and others work as-is with page_pool_iov *

```
static inline dma_addr_t page_pool_get_dma_addr(struct page *page)
{
    dma_addr_t ret;
    if (page_is_page_pool_iov(page))
        return page_pool_iov_dma_addr(page_to_page_pool_iov(page));
    ret = page->dma_addr;
    ...
    return ret;
}
```

Journey of an RX packet: nonpaged memory support

- Drivers must use the page * they receive from page_pool as an opaque token.
- They (almost) already do this!
- page_address() is the main issue.



Journey of an RX packet: incoming packet

- NIC splits the packet into header + payload.
- Payload is DMA'd to a page_pool_iov in device memory.
 - Enables efficient data transfer.
- Header is DMA'd to a header buffer in host memory.
 - Enables the host kernel to parse the packet headers.
- NIC creates a 'devmem' skb and sends it up the stack.

Journey of an RX packet: devmem skb support

- Skbs are required to be either all devmem or host memory.
- Devmem skbs are marked with skb->devmem & skb_frags_not_readable()
- Results in some quirks:
 - Loopback.
 - Software checksum calculation.
 - TCP dump payload access.

Journey of an RX packet: recvmsg() uapi

- Device memory data can't be copied to linear buffer or mapped to user's address space.
- We provide 'pointer' to the memory in the dma-buf accessible from the

userspace.

```
ssize_t ret = recvmsg(client_fd, &msg, MSG_SOCK_DEVMEM);
for (cm = CMSG_FIRSTHDR(&msg); cm; cm = CMSG_NXTHDR(&msg, cm)) {
    if (cm->cmsg_level != SOL_SOCKET ||
        cm->cmsg_type != SCM_DEVMEM_OFFSET){
        continue;
    }
    ...
    cmsg_devmem = (struct cmsg_devmem *)CMSG_DATA(cm);
    ...
    validate_buffer(buf_mem + cmsg_devmem->frag_offset,
        cmsg_devmem->frag_size);
    ...
}
```

Journey of an RX packet: recvmsg() uapi

 Need the user to give us a signal 'setsockopt()' that they're done with the data.

Device memory TX path

- Much more straightforward than RX path.
- Dma-buf to send can be passed to sendmsg() API.
- Largely follows the MSG_ZEROCOPY code path.
- Need to create iov_iter that is backed by page_pool_iovs.
- Re-uses the same devmem skb support as RX path.
- skb_frag_dma_map can grab the dma_addr from page_pool_iovs.

Initial results

- ~96% line rate at TCP level: 192 gbps bi-directional per NIC/GPU pair.
- Running in production.
- Transports exercised with production workloads: NCCL.
- Pytorch exercised with production workloads.
 - Tensorflow, JAX & others use the same transport primitives.
- ~3X better throughput than regular TCP NCCL transports*.
- Comparable network efficiency to RDMA-based NCCL transports for larger message sizes.

Possible follow up work

- io_uring support.
- Dynamic queue management.